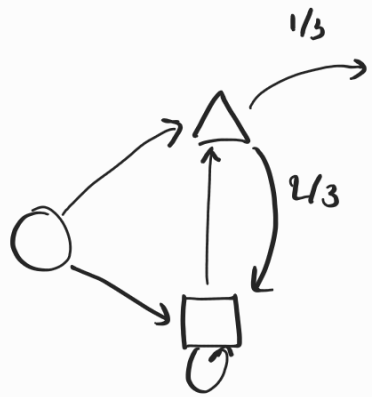


STOCHASTIC GAMES



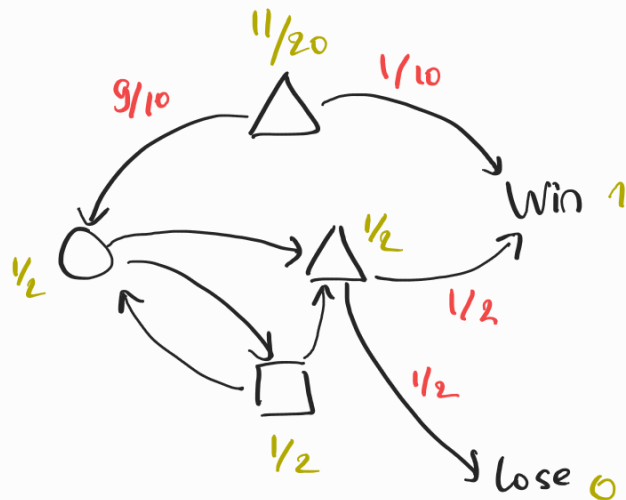
- Max 6
- Min 2
- △ Random

$$V = V_{\text{Max}} (+) V_{\text{Min}} (+) V_{\text{Random}}$$

STOCHASTIC REACHABILITY GAMES

Reach: we fix $\text{Win} \subseteq V$

$$\text{Reach} = \{ \rho = v_0 v_1 \dots \in V^\omega : \exists i v_i \in \text{Win} \}$$



$\text{val} : V \rightarrow [0, 1]$ value function

$$\text{val}^{\sigma, \tau}(v) = \mathbb{P}_v^{\sigma, \tau}(\text{Reach}(\text{Win}))$$

$$= \mathbb{E}_v^{\sigma, \tau}[\mathbb{1}_{\text{Reach}(\text{Win})}]$$

$$\text{val}_{\text{Max}}(v) = \sup_{\sigma} \inf_{\tau} \text{val}^{\sigma, \tau}(v)$$

$$\leq \inf_{\tau} \sup_{\sigma} \text{val}^{\sigma, \tau}(v) = \text{val}_{\text{Min}}(v)$$

Theorem for stochastic reachability games:

- $\text{val}_{\text{Max}} = \text{val}_{\text{Min}} = \text{val}_*$
- pure positional optimal strategies exist

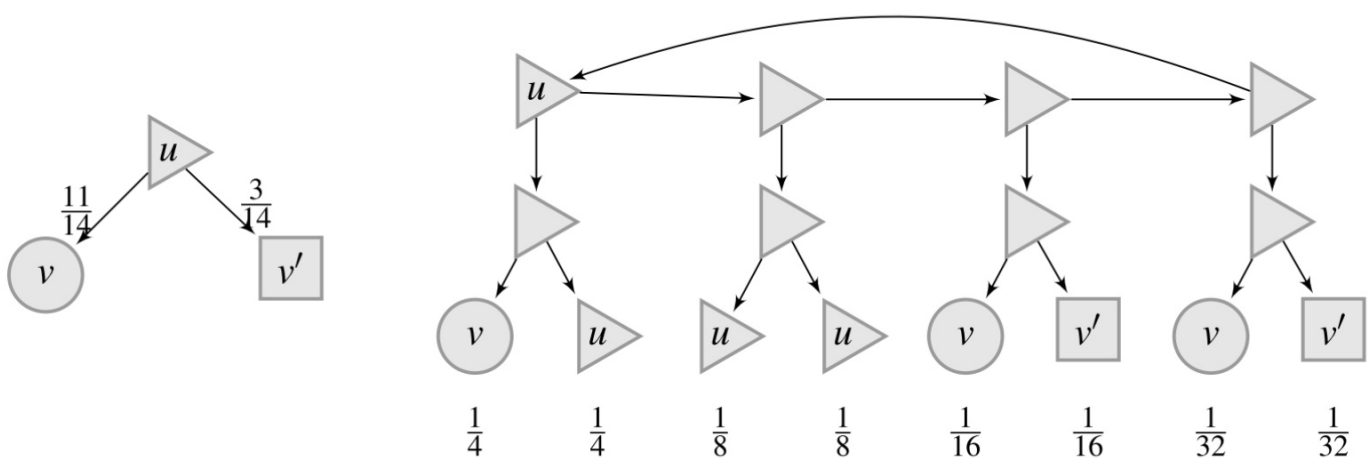
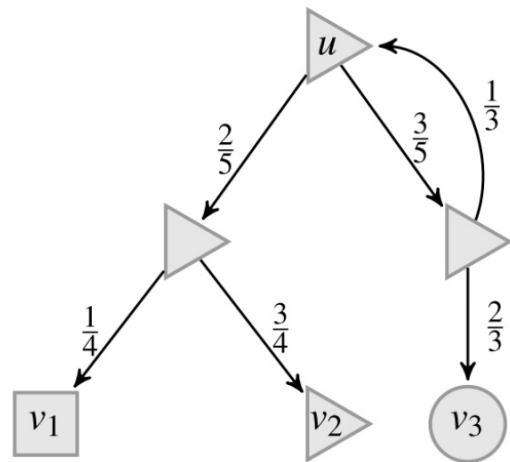
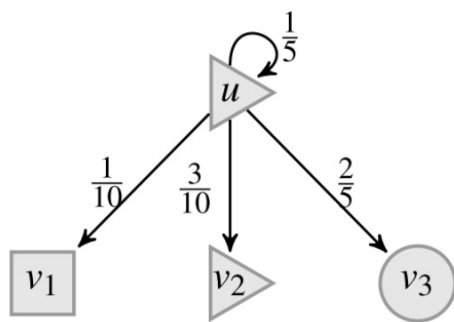
$$\begin{cases} \sigma : V \xrightarrow{\text{Max}} E \\ \tau : V \xrightarrow{\text{Min}} E \end{cases}$$

SIMPLE STOCHASTIC GAMES

stochastic reachability games where:

- outdegree = 2
- probabilities are binary: $\frac{1}{2}/\frac{1}{2}$

REDUCTION



Theorem (Knaster-Tarski)

(\mathcal{L}, \leq) complete lattice

$$\phi: (\mathcal{L}, \leq) \rightarrow (\mathcal{L}, \leq)$$

monotonic $x \leq y \Rightarrow \phi(x) \leq \phi(y)$

preserves suprema: $\sup_n \phi(x_n) = \phi(\sup_n x_n)$

(1) ϕ has a unique least fixpoint: $\phi(x) = x$

(2) the least fixpoint is also the least prefixpoint: $\phi(x) \leq x$

(3) $x_0 = \perp$ $x_{k+1} = \phi(x_k)$

$\lim_k x_k$ is the least fixed point of ϕ

$$Y = \{ \mu: V \rightarrow [0,1] \}$$

$$\mu \leq \mu' \text{ if } \forall v \in V \mu(v) \leq \mu'(v)$$

(Y, \leq) is a complete lattice

$$\textcircled{1} : Y \rightarrow Y$$

$$\textcircled{1}(\mu)(u) = \begin{cases} 1 & \text{if } u \in W_{\text{in}} \\ \max \{ \mu(v) : u \rightarrow v \} & u \in V_{\text{max}} \\ \min \{ \mu(v) : u \rightarrow v \} & u \in V_{\text{min}} \\ \sum_v \delta(v|u) \mu(v) & u \in V_{\text{Random}} \end{cases}$$

$\textcircled{1}$ monotonic, preserves suprema

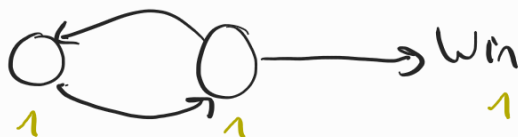
Theorem: val^* = Least Fixed Point of $\textcircled{1}$

if $\forall u \in V_{\text{min}}$

$z(u) \in \text{argmin} \{ \text{val}^*(v) : u \rightarrow v \}$

then z is optimal

⚠ False for Max:



$$\mu_0 = 0 \quad \mu_{k+1} = \mathbb{D}(\mu_k)$$

$$\mu_* = \lim_k \mu_k \quad \mu_* = \text{LFP}(\mathbb{D})$$

$$\text{Reach}_{\leq k} = \{ \rho = v_0 v_1 \dots \in V^\omega : \exists i \leq k \ v_i \in W_{\text{in}} \}$$

Claims

$$(1) \quad \mu_k = \sup_{\delta} \inf_{\gamma} P^{\delta, \gamma}(\text{Reach}_{\leq k})$$

$$(2) \quad \mu_* \leq \text{val}_{\max}$$

[val_{\max} is a prefixed point of \mathbb{D}]

$$(3) \quad \text{val}_{\min} \leq \mu_*$$

[invariant : $\forall \delta \forall k \forall u$
 $\mathbb{E}_{u, \delta} [\mu^*(\pi_k)] \leq \mu^*(u)$]

POSITIONAL DETERMINACY FOR MAX

Fix $\lambda \in (0, 1]$

$$\text{QReach}_\lambda(Q) = \begin{cases} 0 & \text{if } Q \notin \text{Reach} \\ \lambda^i & \text{if } i = \text{first } v_i \in \text{Win} \end{cases}$$

$$Y = \{ \mu : V \rightarrow [0, 1] \}$$

$$\mathbb{D}_\lambda : Y \rightarrow Y$$

$$\mathbb{D}_\lambda(\mu)(u) = \begin{cases} 1 & \text{if } u \in \text{Win} \\ \lambda \max \{ \mu(v) : u \rightarrow v \} & u \in V_{\max} \\ \lambda \min \{ \mu(v) : u \rightarrow v \} & u \in V_{\min} \\ \lambda \sum_v \delta(v|u) \mu(v) & \end{cases}$$

\mathbb{D}_λ monotonic, preserves suprema, and contracting

Theorem (Banach)

$(X, \|\cdot\|)$ complete space

$$\Phi: X \rightarrow X$$

λ -contracting: $\|\Phi(x) - \Phi(y)\| \leq \lambda \|x - y\|$

Φ has a unique fixed point x_*

$$\forall x_0 \quad \lim_k \Phi^k(x_0) = x_*$$

$$\|\Phi^k(x_0) - x_*\| \leq \frac{\lambda^k}{1-\lambda} \|\Phi(x_0) - x_0\|$$

Step 1: study QReach_λ

σ optimal $\Leftrightarrow \forall u \in V_{\max} \quad \sigma(u) \in \text{argmax} \{ \text{val}_v(v) : u \rightarrow v \}$

τ optimal $\Leftrightarrow \forall u \in V_{\min} \quad \sigma(u) \in \text{argmin} \{ \text{val}_v(v) : u \rightarrow v \}$

Step 2: $\lim_{\lambda \rightarrow 1} \text{QReach}_\lambda = \text{Reach}$

$\exists (\lambda_k)_k \rightarrow 1$ all use the same σ pure positional

Claim: σ optimal for Reach

(1) $\forall \mu: V \rightarrow [0, 1] \quad \forall \lambda$

$$\| \mathbb{D}_\lambda(\mu) - \mathbb{D}(\mu) \| \leq 1 - \lambda$$

(2) val_{λ_k} fixed point of \mathbb{D}_{λ_k}

$$\text{So } \| \text{val}_{\lambda_k} - \mathbb{D}(\text{val}_{\lambda_k}) \| \leq 1 - \lambda_k$$

(3) $\lim_k \text{val}_{\lambda_k}$ is a fixed point of \mathbb{D}

$$\text{So } \lim_k \text{val}_{\lambda_k} \geq \text{val}_*$$

(4) $\mathbb{P}_u^{\sigma, \tau}(\text{Reach}) \geq \mathbb{P}_u^{\sigma, \tau}(\text{QReach}_{\lambda_k}) \geq \text{val}_{\lambda_k}(u)$

$k \rightarrow \infty$

$$\forall u \quad \mathbb{P}_u^{\sigma, \tau}(\text{Reach}) \geq \lim_k \text{val}_{\lambda_k}(u) = \text{val}_*(u)$$

So: σ is optimal for Reach